

# Relational-Style XML Query

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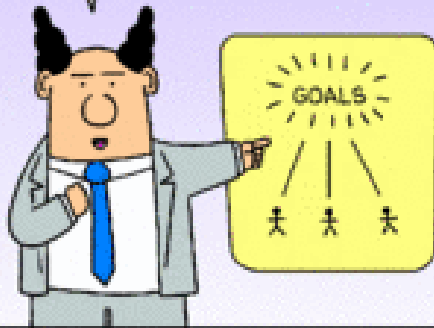
University of Tokyo

June 10<sup>th</sup>, SIGMOD 2008  
Vancouver, Canada

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# If Your Manager Says ...

I DECIDED TO  
START A NEW  
XML PROJECT

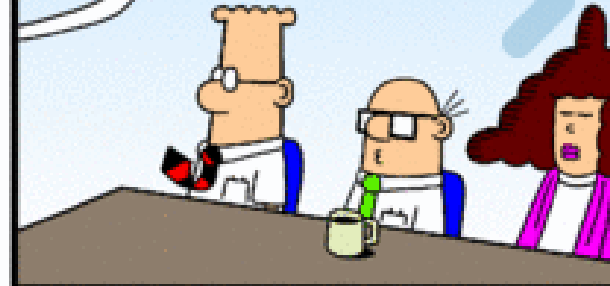


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MASTERING XML IS  
CRUCIAL TO OUR  
COMPANY BECAUSE IT  
IS COMPLETELY A  
NEW DATA MODEL.



EVERYBODY MUST START  
LEARNING SAX, DOM,  
XPath, XQuery, DTD,  
XML Schema, Relax NG...



It's a kind of tragedy...

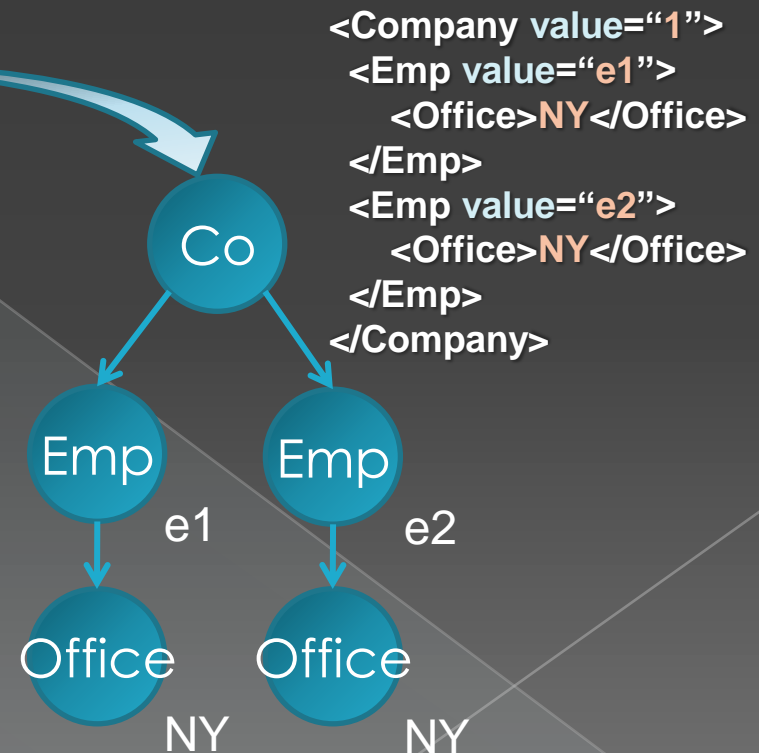
# Migration to XML Database

- Benefits of using XML:

- > XML is a portable text-data format
- > Tree-structured XML can reduce redundancy of relational data.

Company	Employee	Office
1	e1	NY
1	e2	NY

Relational Data



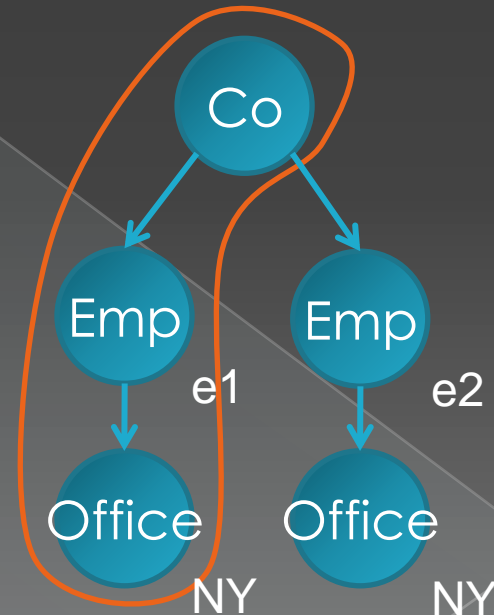
XML Data

# Problem

- Querying relational data translated into XML
- Q: Retrieve a node tuple (Co, Emp, Office) from the XML data
  - > e.g. XPath, a path expression query /Co/Emp/Office

Co	Emp	Office
1	e1	NY
1	e2	NY

**Relational Data**



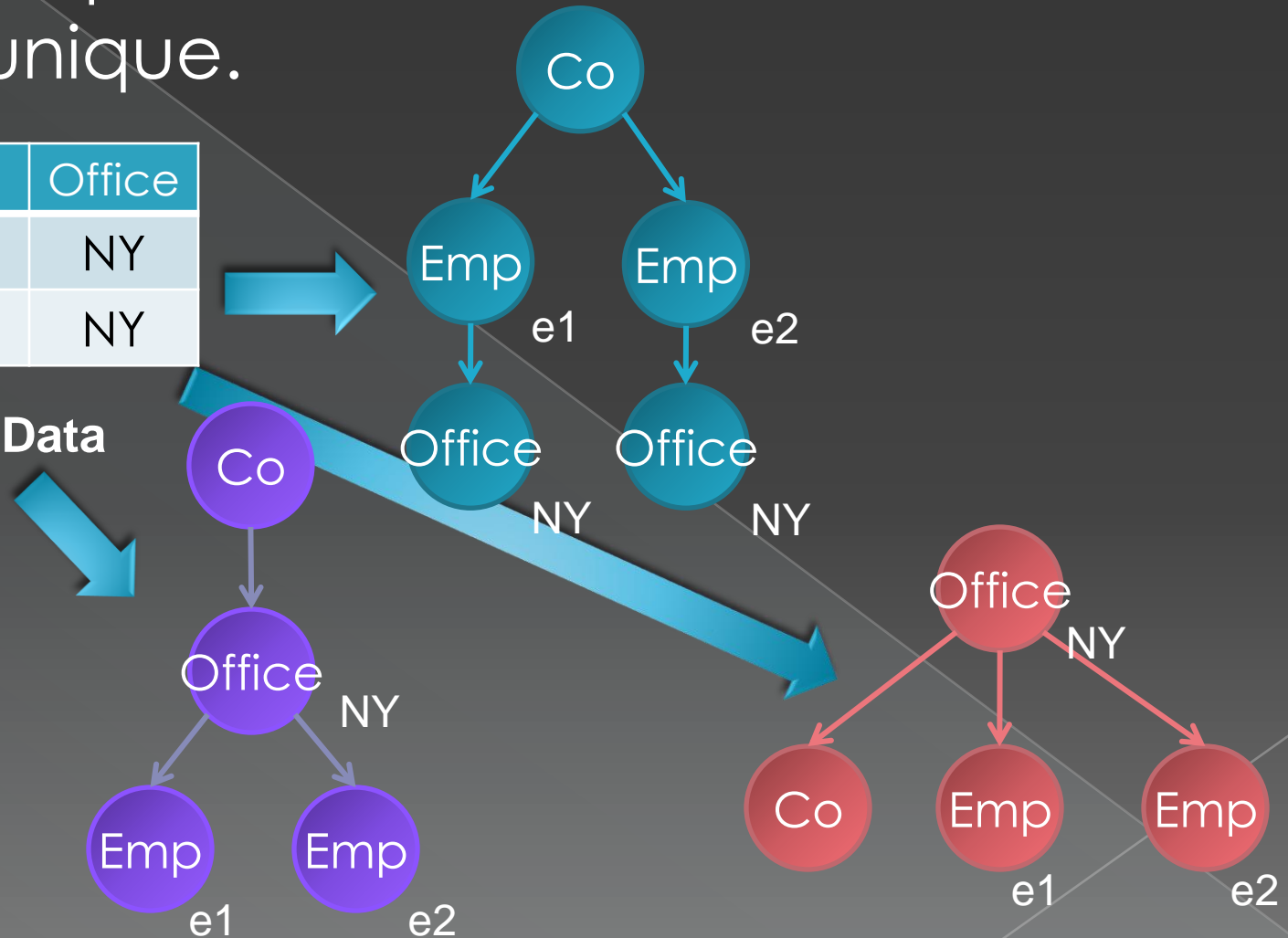
**XML Data**

# A Pitfall: Structural Variations

- Tree-representation of relational data is not unique.

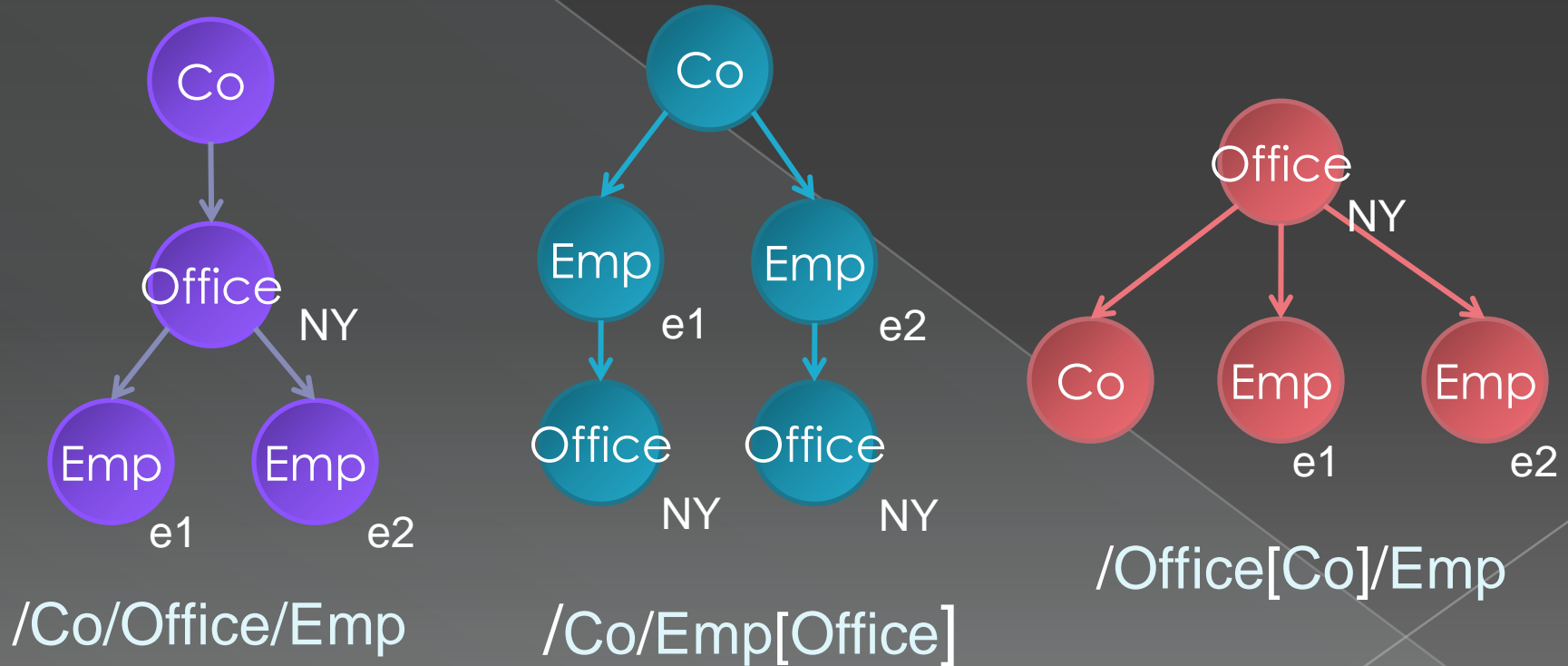
Co	Emp	Office
1	e1	NY
1	e2	NY

**Relational Data**



# Inconvenience of XPath Query

- User must know the entire XML structures to produce correct path queries.



[X] : twig node to test

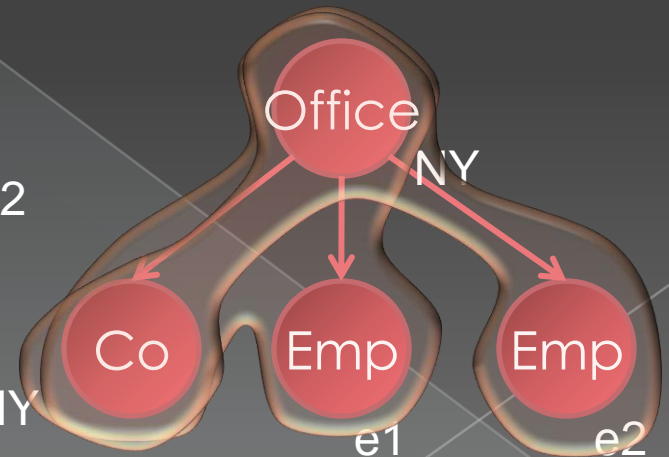
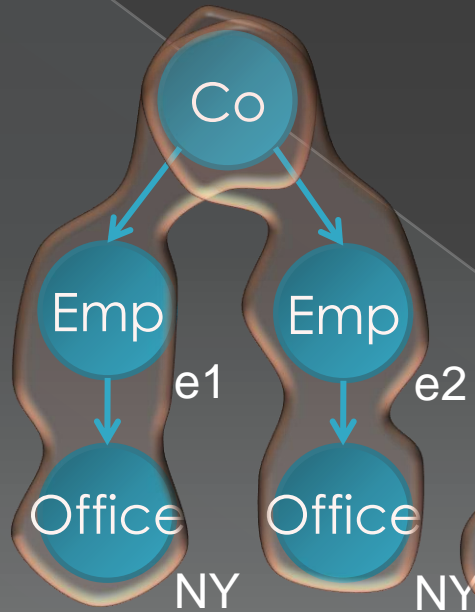
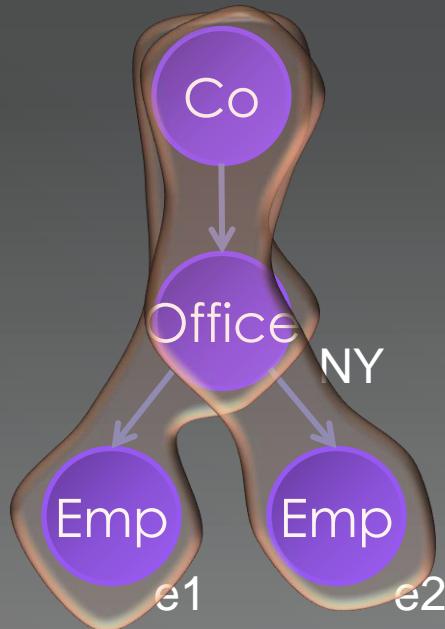
# Relation in XML

- A key observation:

- Relation is simply embedded in XML

Co	Emp	Office
1	e1	NY
1	e2	NY

**Relational Data**



# To Retrieve Relations in XML...

WHY DO WE HAVE TO USE XPATH?

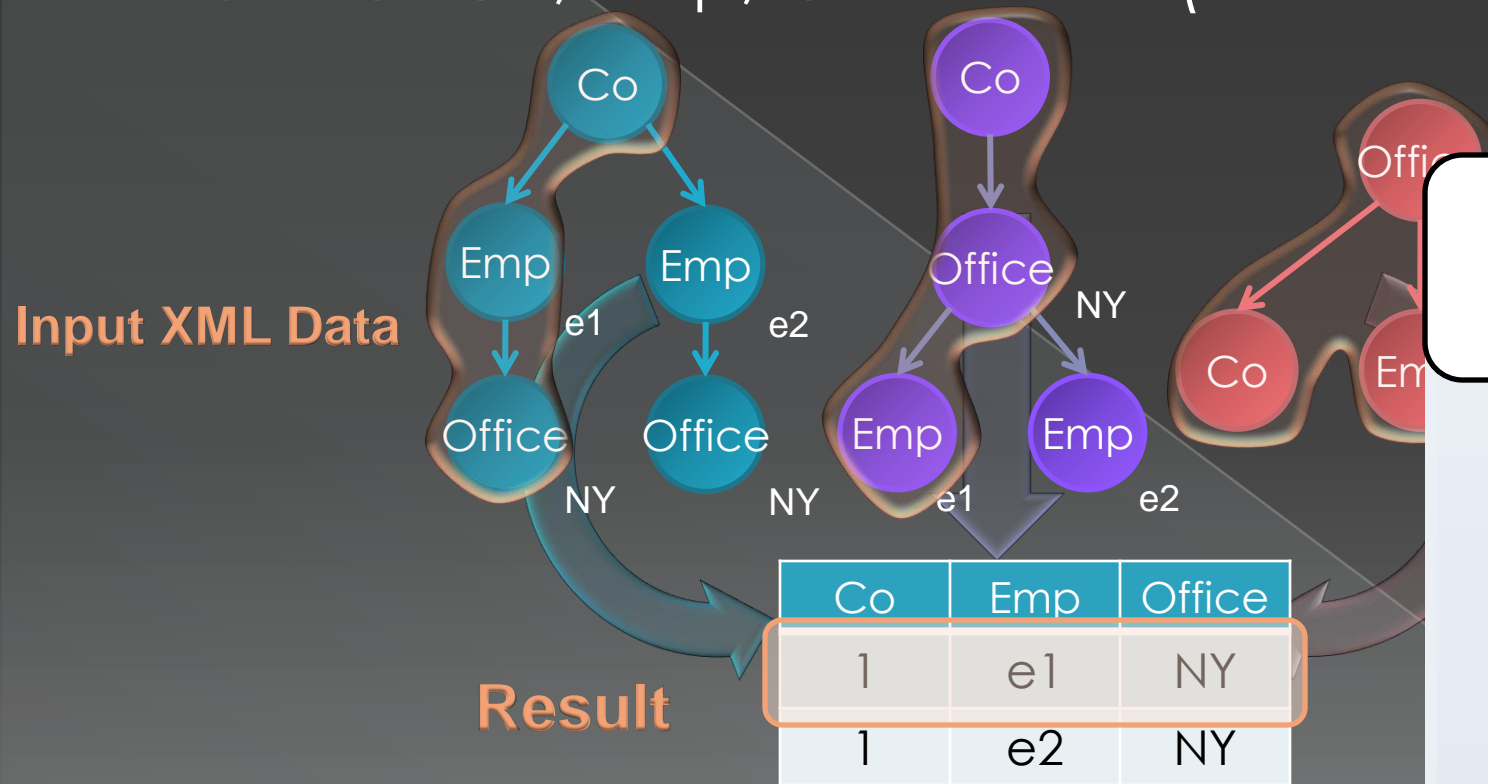


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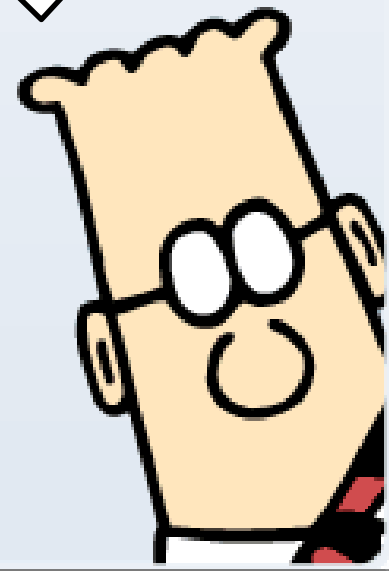
# Relational-Style XML Query

- Query relations in XML
  - > with an SQL-like syntax
- SELECT** Co, Emp, Office **from** (XML Data)



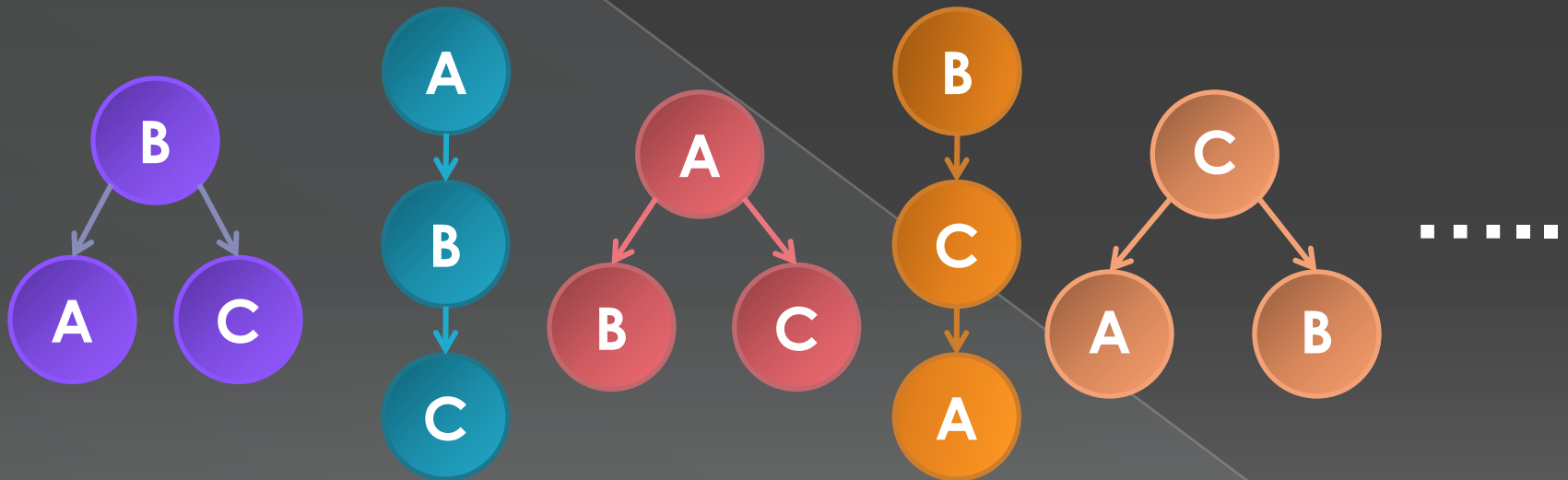
SQL over XML!

- The query statement is stable for variously structured XML data



# Problem Definition

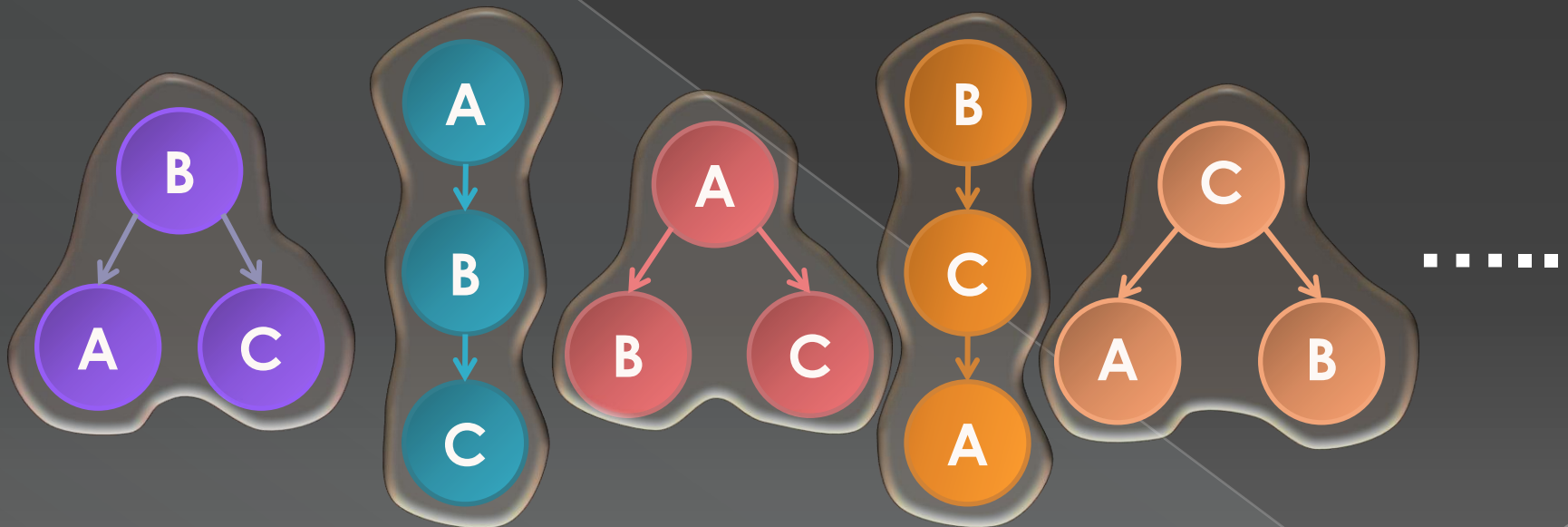
- Convert an SQL query, **SELECT A, B, C**, into an XML structure query.
  - There can be many structural variations of (A, B, C)



- For  $N$  nodes, there exists  $N^{N-1}$  structural variations.

# Amoeba

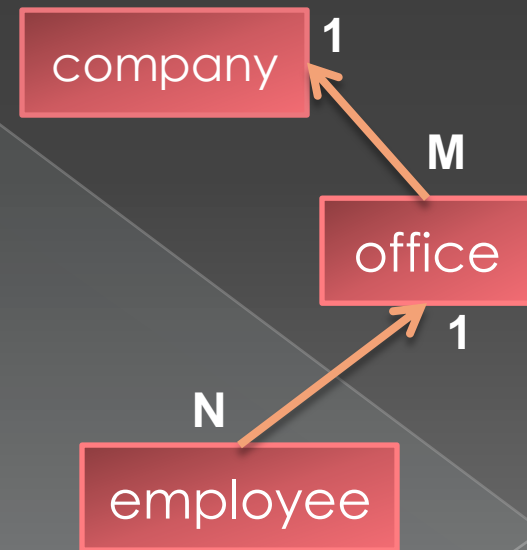
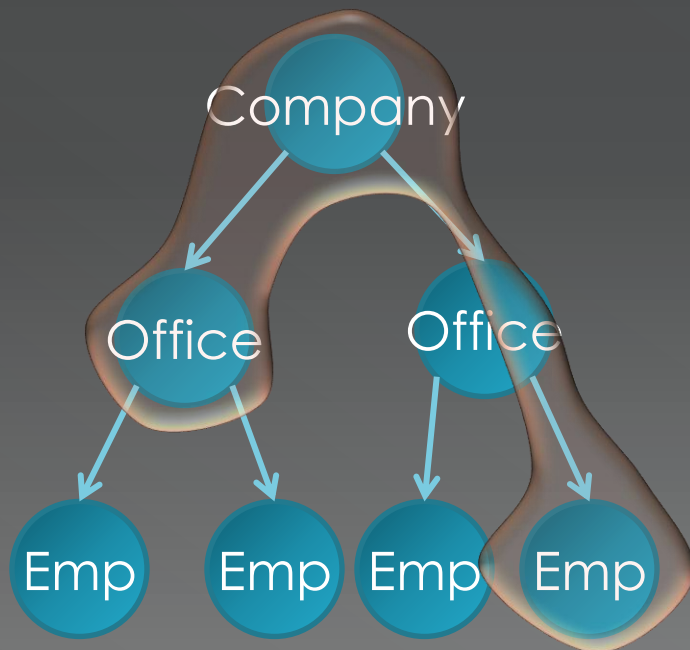
- A node tuple (A, B, C) is an **amoeba** iff one of the A, B and C is a common ancestor of the others.



- **Amoeba join** retrieves all amoeba structures in the XML data.

# Hidden Semantics in XML

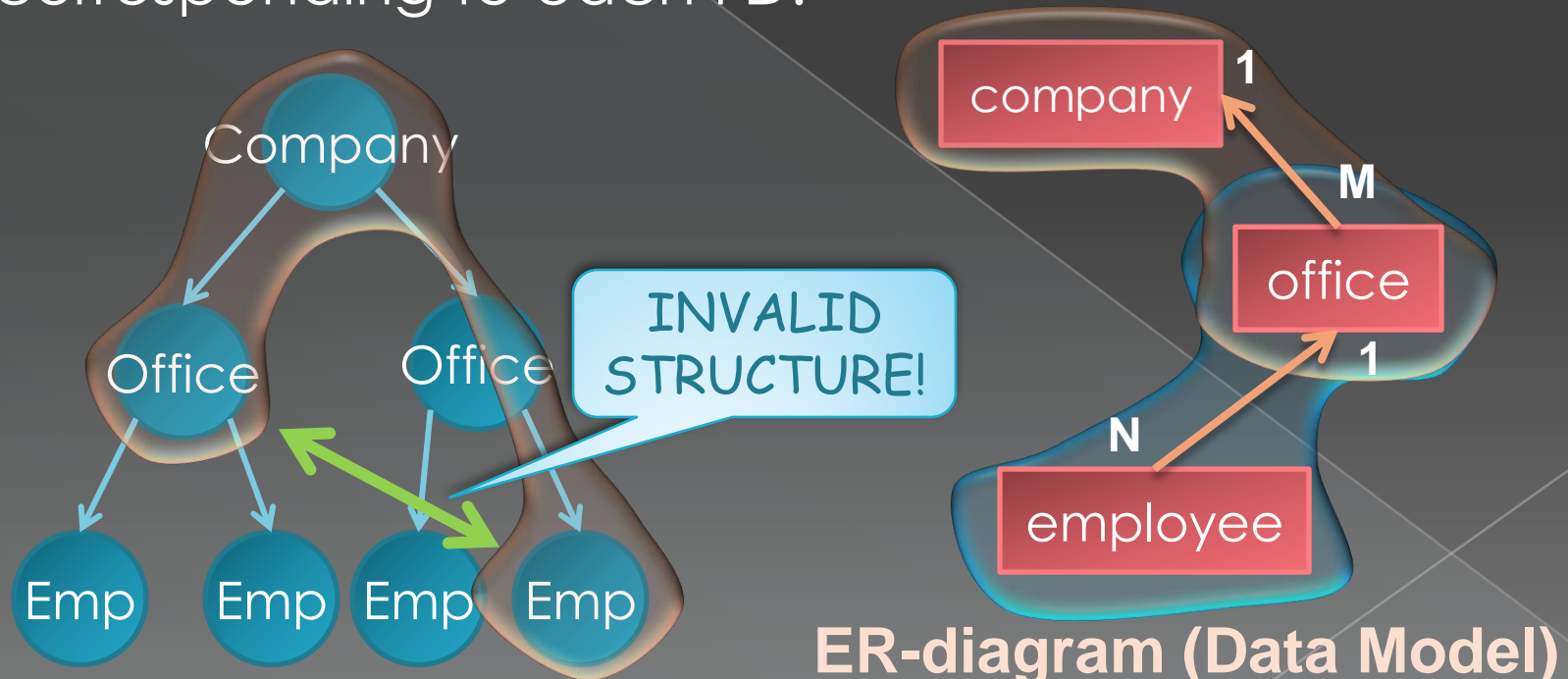
- Some amoeba structure may not form a relation.
  - > Why this structure is not allowed?
- Because there are **functional dependencies (FD)** implied in the XML structure.



ER-diagram (Data Model)

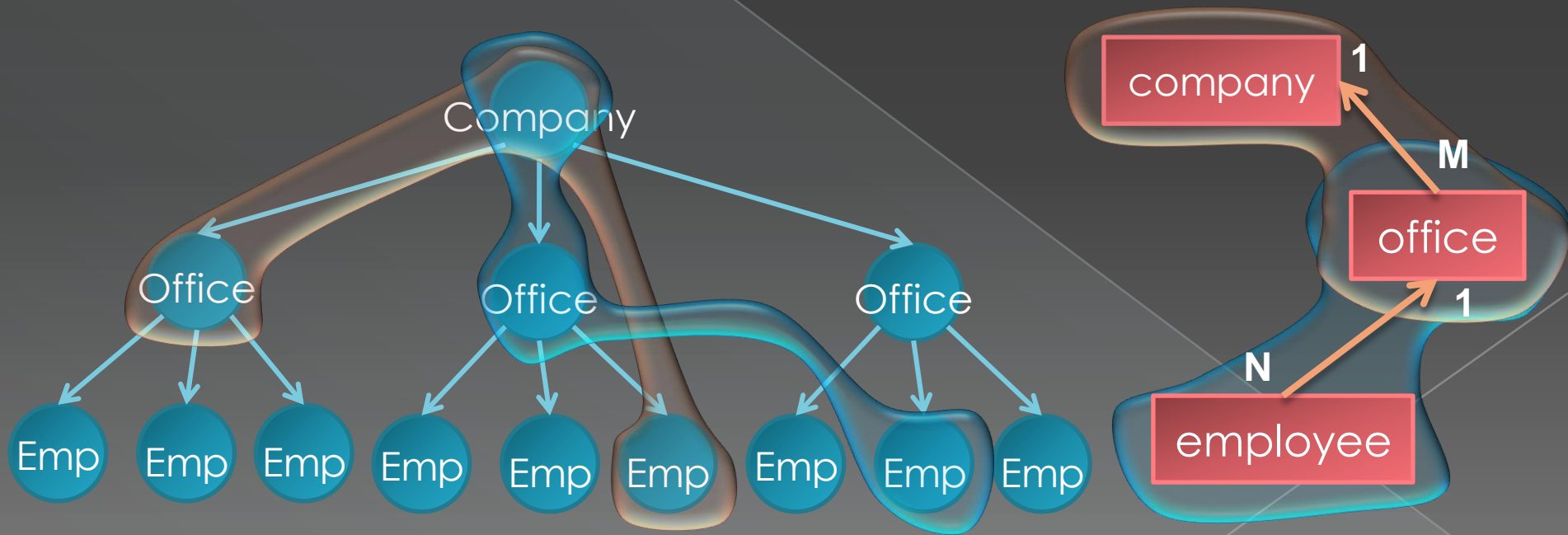
# Functional Dependencies (FD)

- FD:  $X \rightarrow Y$  (From a given  $X$ ,  $Y$  is uniquely determined)
  - > **employee  $\rightarrow$  office** (Each employee belongs to an office)
  - > **office  $\rightarrow$  company** (Each office belongs to a company)
- Relation in XML must have an amoeba structure corresponding to each FD.



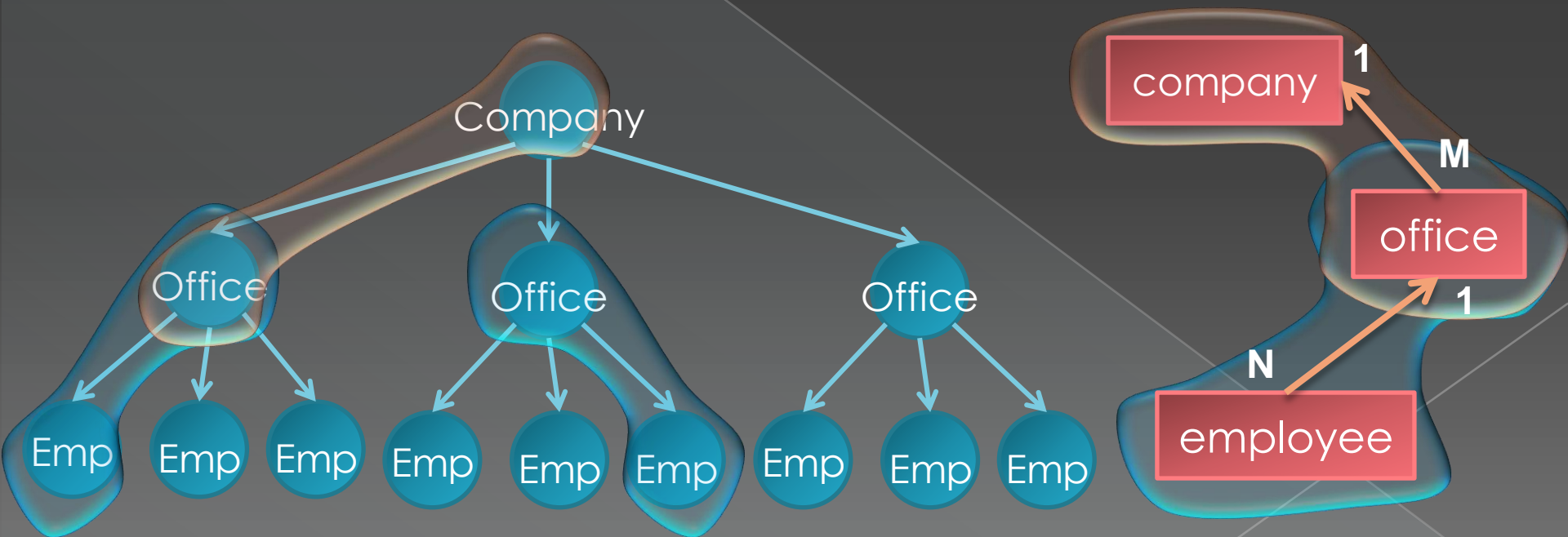
# If FDs are ignored....

- The company has  $M$  offices, and each office has  $N$  employees:
- # of (company, office, employee) tuples:
  - > When  $M = 100, N = 5$   $100 \times (100 \times 5) = 50,000$
- While, # of correct answers is only  $M * N = 500$



# FD-Aware Amoeba Join

- FDs:  $\text{Emp} \rightarrow \text{Office}$ ,  $\text{Office} \rightarrow \text{Company}$
- Bottom-up construction of query results
  - Amoeba Join (Employee, Office)
  - Amoeba Join (Office, Company)

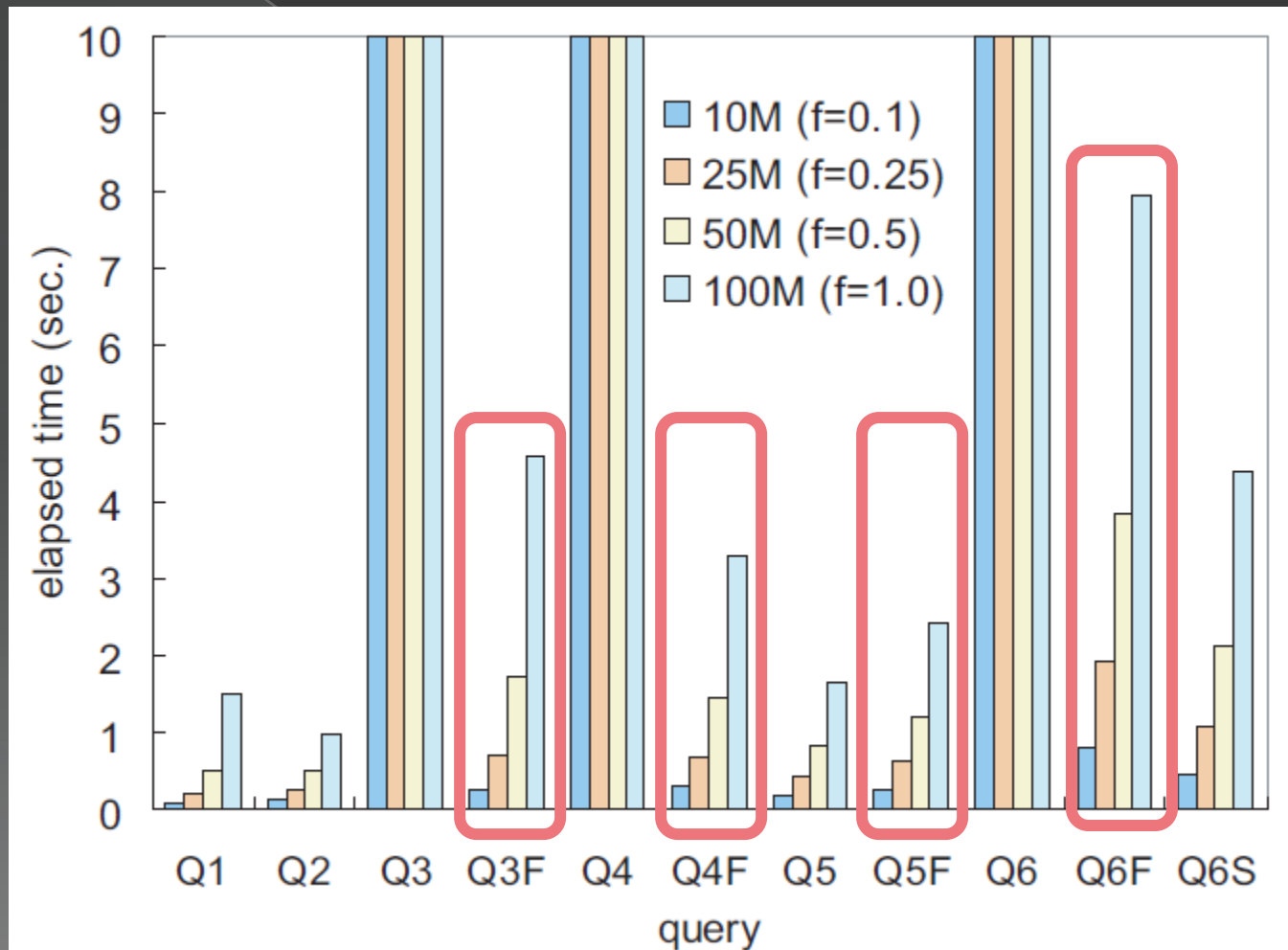


- FD-aware amoeba join avoids invalid XML structures.



# Query Performance

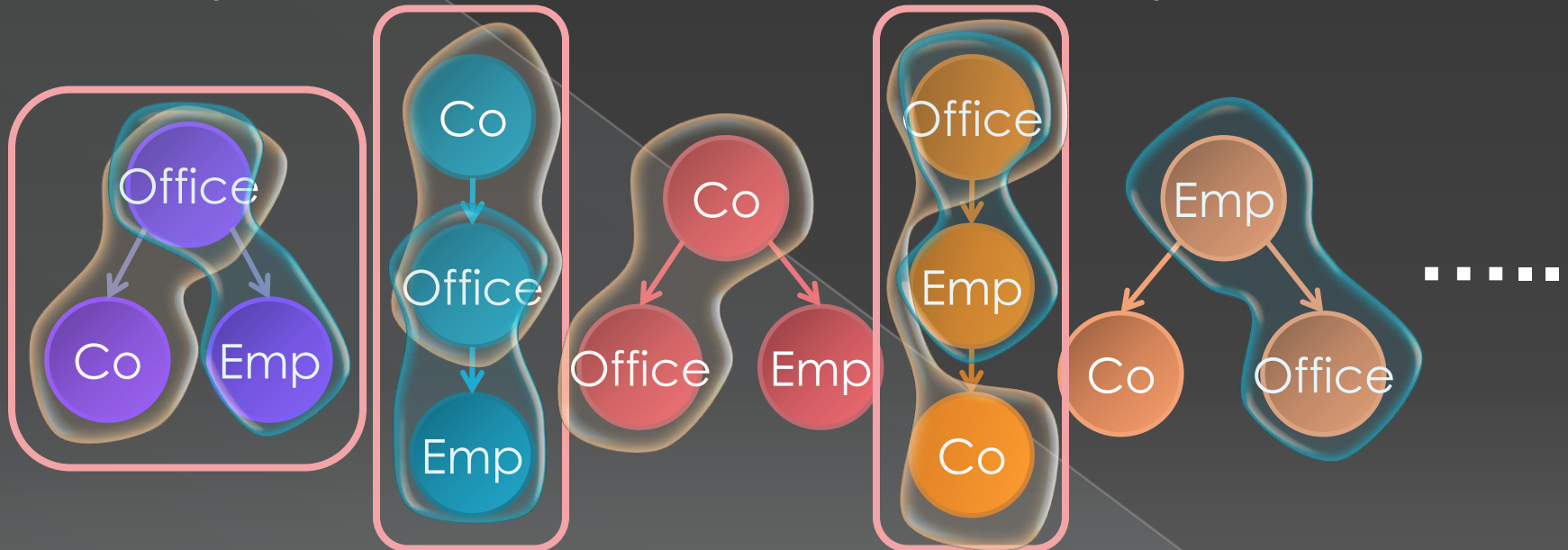
- FD-aware amoeba join scales well
  - For various sizes of XML data





# Query Translation using FDs

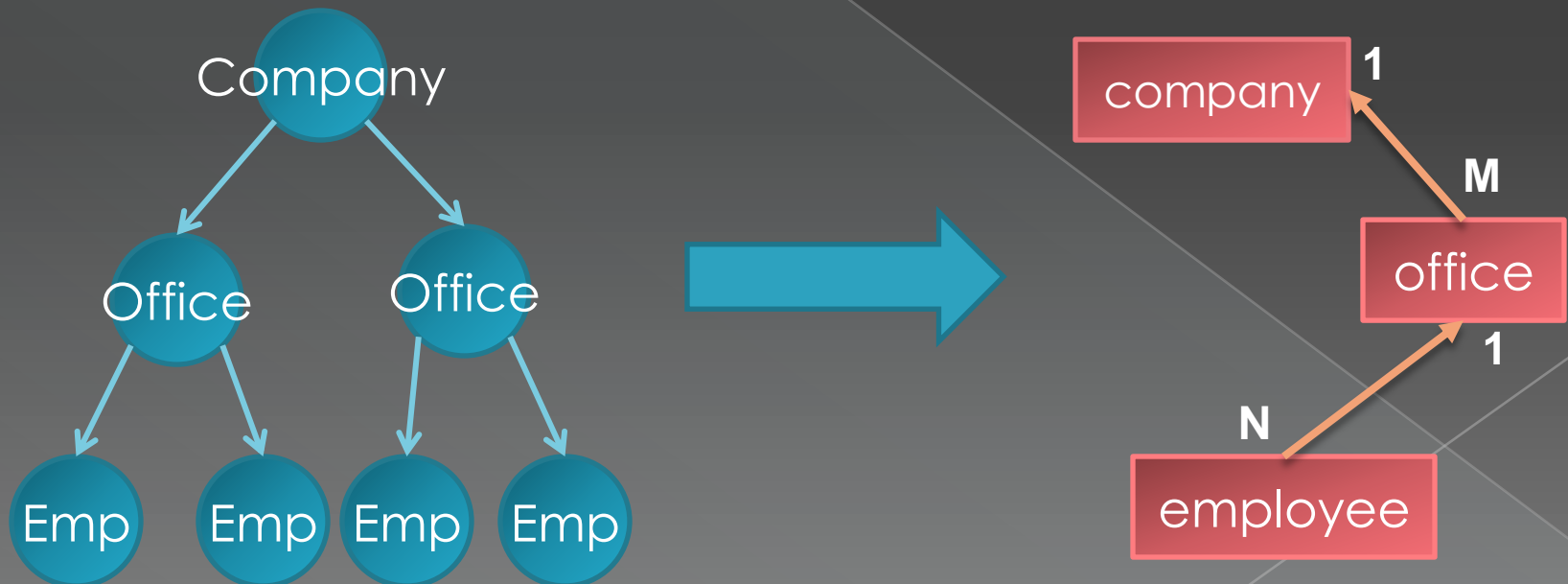
- Relational query into XML query
  - SELECT Co, Office, Emp
    - (with FDs:  $\text{Emp} \rightarrow \text{Office}$ ,  $\text{Office} \rightarrow \text{Co}$ )



- XML structures of interest are automatically determined from a relation and functional dependencies

# Detecting FDs

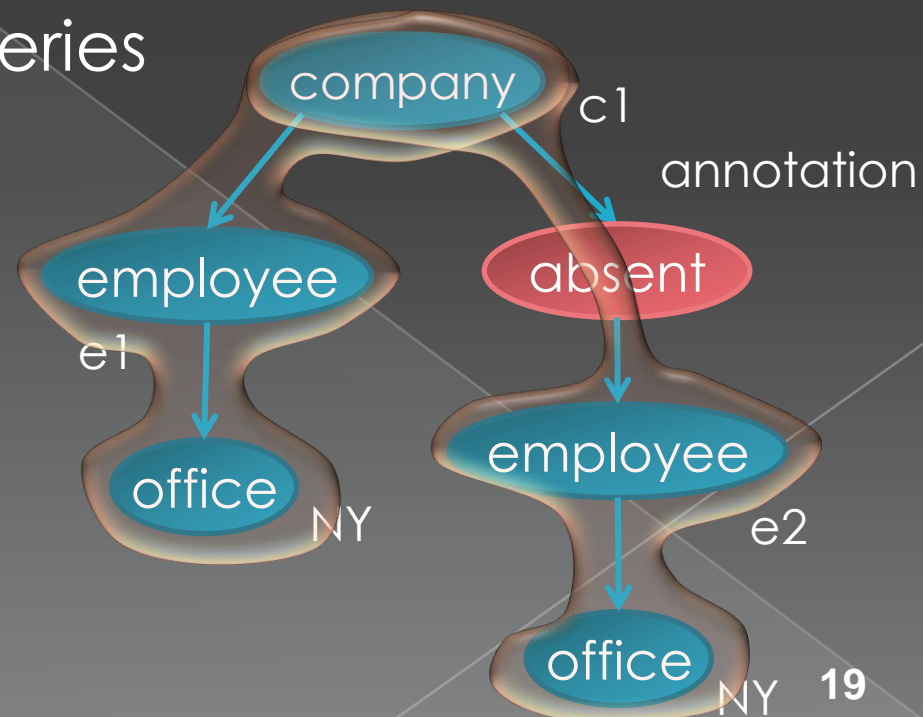
- A type of FDs required to determine XML structures to query is **one-to-many (or one-to-one)** relationships:
  - FD: Emp  $\rightarrow$  Office
    - Each employee belongs to an office
    - An office may have several employees (one-to-many)
- We can observe these relationships by counting node occurrences or directory from the ER-diagram.



# Think in Relational-Style

- First, consider
  - > XML := Relations + their annotations
- Steps
  - > 1. Detect relational part from XML data
  - > 2. Detect one-to-many(one) relationships (FDs)
  - > 3. Write relational queries
    - SELECT Co, Emp, Office

- Note:
  - It is also possible to include annotations in query statements.



# Summary of Our Contributions

- Relation in XML

- > Defined using amoeba structure and FDs

- Relational-Style XML Query

- > Retrieves relations in XML with a SQL-like query syntax (SQL over XML)
  - > Allows structural variations of XML data

- Departure from path expression queries

- > Target XML structures are automatically determined.

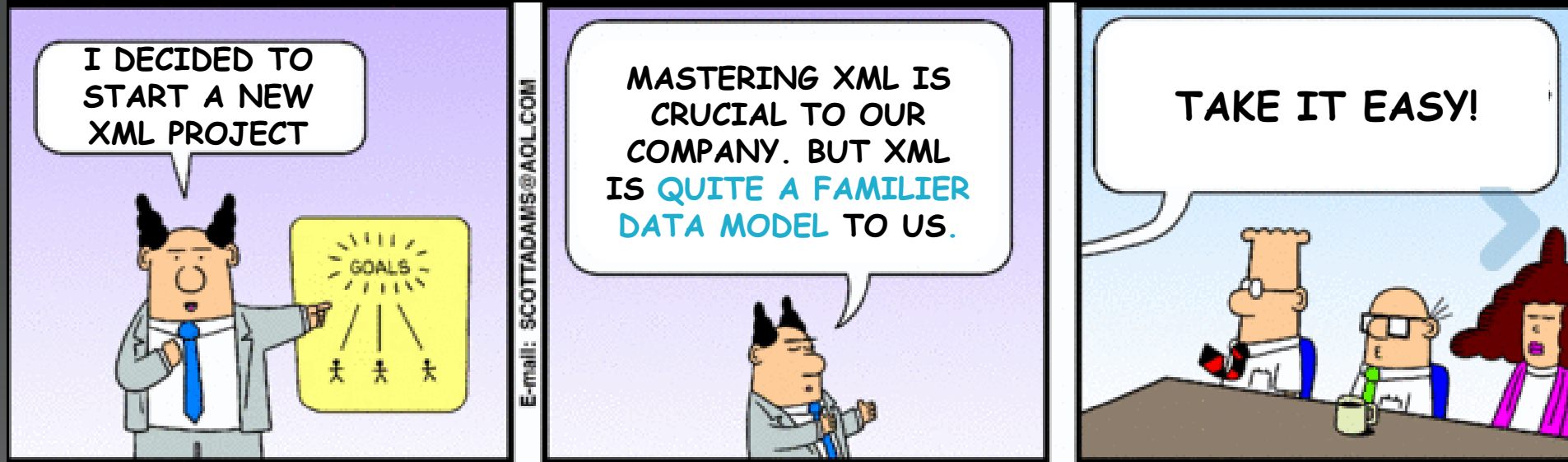
# Applications of Relational-Style

- ◉ (see the paper for details)
- ◉ XML Algebra
  - > Based on relational-semantics
    - selection, projection, etc.
- ◉ Keys for XML
  - > A key is a special-case of FDs
- ◉ Database integration
- ◉ Schema evolution
- ◉ Managing relational data enhanced with XML syntax
- ◉ A lot more...

# Conclusions

- “It’s Just SQL”

- A large number of XML data and queries are still relational.



- Before going deep into the XML world,  
**Think in Relational-Style!!!**